

The Zero-Free Intervals for Characteristic Polynomials of Matroids

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Let M be a loopless matroid with rank r and c components. Let $P(M, t)$ be the characteristic polynomial of M . We shall show that $(-1)^r P(M, t) \geq (1-t)^r$ for $t \in (-\infty, 1)$, that the multiplicity of the zeros of $P(M, t)$ at $t = 1$ is equal to c , and that $(-1)^{r+c} P(M, t) \geq (t-1)^r$ for $t \in (1, \frac{32}{27}]$. Using a result of C. Thomassen we deduce that the maximal zero-free intervals for characteristic polynomials of loopless matroids are precisely $(-\infty, 1)$ and $(1, \frac{32}{27}]$.

1. Introduction

Characteristic polynomials of matroids were first studied by Rota [6]. Heron [3] defined chromatic polynomials of matroids and showed that they are equivalent to characteristic polynomials. Here, we are concerned with arbitrary matroids, so we will refer to *characteristic polynomials*, reserving the term *chromatic polynomials* for graphs. It is known (see Jackson [4], Tutte [9] and Woodall [12]) that the chromatic polynomial of an arbitrary graph $P(G, t)$ has no zeros in the intervals $(-\infty, 0)$, $(0, 1)$ and $(1, \frac{32}{27}]$, that the multiplicity of zeros at 0 is equal to the number of components of the graph and that the multiplicity of zeros at 1 is equal to the number of blocks of the graph. Recently Thomassen [8] has shown that the chromatic zeros of graphs are dense in $(\frac{32}{27}, \infty)$. It follows that the maximal zero-free intervals for chromatic polynomials of graphs are precisely $(-\infty, 0)$, $(0, 1)$ and $(1, \frac{32}{27}]$. He also showed that the chromatic polynomial $P(G, t)$ of a connected graph G with n vertices satisfies $|P(G, t)| \geq |t(t-1)^{n-1}|$ for all $t \leq 32/27$.

The facts that the characteristic polynomial of a graphic matroid has no zeros in the intervals $(-\infty, 1)$ and $(1, \frac{32}{27}]$ and that the multiplicity of zeros at 1 is equal to the number of connected components of the matroid follow from the above results on chromatic polynomials of graphs and the well-known equality $t^c P(M, t) = P(G, t)$ between the chromatic polynomial $P(G, t)$ of a graph G with c components, and the characteristic polynomial $P(M, t)$ of its associated graphic matroid M . It is also known that the

characteristic polynomial of a cographic matroid has no zeros in the intervals $(-\infty, 1)$ and $(1, \frac{32}{27}]$, and that the multiplicity of zeros at 1 is equal to the number of connected components of the matroid. This follows immediately from the corresponding results for flow polynomials of graphs (see Wakelin [11]). We shall deduce that both these zero-free intervals hold for the characteristic polynomial $P(M, t)$ of a loopless matroid M with c components and rank r by showing that $(-1)^r P(M, t) \geq (1-t)^r$ for $t \in (-\infty, 1)$ and $(-1)^{r+c} P(M, t) \geq (t-1)^r$ for $t \in (1, \frac{32}{27}]$. We will also show that the multiplicity of the zeros of $P(M, t)$ at $t = 1$ is equal to c .

Given a matroid M , we shall use S_M , \mathcal{C}_M , \mathcal{I}_M and r_M to denote the groundset, the set of circuits, the set of independent sets and the rank function, respectively, of M . These will be simplified to S , \mathcal{C} , \mathcal{I} and r when it is obvious which matroid we are referring to. We denote the rank of M by $r(M)$, the number of elements of M by $|M|$ and the number of components of M by $c(M)$. Thus $r(M) = r_M(S)$ and $|M| = |S|$. Let $M - e$ and M/e respectively denote the matroids resulting from the deletion and contraction of an element e .

Sections 2–5 contain definitions and lemmas that will be used in the proofs of the main results of the paper in Section 6. We refer the reader to Oxley [5] for basic definitions not given in this paper.

2. Some elementary matroid results

Lemma 2.1. ([2], Proposition 4.1) *Let M be a matroid with $e \in S$, $A \subseteq S - e$. Then*

- (a) $r_{M-e}(A) = r_M(A)$
- (b) $r(M - e) = r(M) - 1$ if e is a coloop of M
- (c) $r(M - e) = r(M)$ if e is not a coloop of M .

By repeated application of Lemma 2.1(a) we obtain the following lemma.

Lemma 2.2. *Let M be a matroid with $A, B \subseteq S$, $A \cap B = \emptyset$. Then $r_M(A) = r_{M-B}(A)$.*

Lemma 2.3. ([2], Proposition 4.1) *Let M be a matroid with $e \in S$, $A \subseteq S - e$. Suppose e is not a loop in M . Then*

- (a) $r_{M/e}(A) = r_M(A \cup \{e\}) - 1$
- (b) $r(M/e) = r(M) - 1$.

Lemma 2.4. ([5], Proposition 3.1.26) *Let M be a matroid with $e, f \in S$. Then $(M - e)/f = (M/f) - e$.*

Lemma 2.5. ([9], Theorem 6.5) *Let $e \in S$ be an element of a connected matroid M . Then either $M - e$ or M/e is connected.*

Lemma 2.6. ([9], Theorem 6.6) *Let e be an element of a 3-connected matroid M . Then, provided $|M| \geq 4$, both $M - e$ and M/e are connected.*

3. Characteristic polynomials

The characteristic polynomial $P(M, t)$ of a matroid M is a polynomial in t defined by

$$P(M, t) = \sum_{A \subseteq S} (-1)^{|A|} t^{r_M(S) - r_M(A)}.$$

Lemma 3.1. ([3], Lemma 1.4) *If the matroid M has a loop, then $P(M, t) \equiv 0$.*

Lemma 3.2. ([3], Lemma 1.2) *If M is a matroid and $e \in S$ is not a loop or a coloop, then*

$$P(M, t) = P(M - e, t) - P(M/e, t).$$

Lemma 3.3. ([3], Lemma 1.5) *If e is a coloop in a loopless matroid M , then $P(M, t) = (t - 1)P(M - e, t)$.*

Lemma 3.4. ([3], Lemma 3.10) *Let M be a matroid and M_1, M_2, \dots, M_n be the components of M . Then $P(M, t) = \prod_{i=1}^n P(M_i, t)$.*

Lemma 3.5. *If e is a parallel element in a matroid M , then $P(M, t) = P(M - e, t)$.*

Proof. Follows from Lemmas 3.1 and 3.2. □

4. Parallel connections, series connections and 2-sums

Let M_1 and M_2 be matroids. We shall denote the direct sum of M_1 and M_2 by $M_1 \oplus M_2$. We shall denote the parallel connection, series connection, and 2-sum of M_1 and M_2 about basepoint p by $M_1 \parallel^p M_2$, $M_1 *^p M_2$, and $M_1 \oplus_2^p M_2$, respectively.

Given matroids M_i , $1 \leq i \leq n$, such that $S_{M_i} \cap S_{M_j} = \{p\}$ for $1 \leq i < j \leq n$, where p is not a loop or a coloop in M_i and $|M_i| \geq 3$ for $1 \leq i \leq n$, we define the 2-sum of M_1, \dots, M_n to be the matroid

$$M_1 \oplus_2^p M_2 \oplus_2^p \cdots \oplus_2^p M_n = (M_1 \parallel^p M_2 \parallel^p \cdots \parallel^p M_n) - p.$$

Lemma 4.1. ([5], Corollary 7.1.23) *The class of graphic matroids is closed under parallel connection, series connection and 2-sum.*

Lemma 4.2. *Let $M = M_1 \parallel^p M_2 \parallel^p \cdots \parallel^p M_n$.*

- (a) $r(M) = \sum_{i=1}^n r(M_i) - n + 1$.
- (b) $c(M) = \sum_{i=1}^n c(M_i) - n + 1$.
- (c) $M/p = M_1/p \oplus M_2/p \oplus \cdots \oplus M_n/p$.
- (d) $P(M, t) = (t - 1)^{-n+1} \prod_{i=1}^n P(M_i, t)$.

Proof. Follows from [1, Theorem 6.16(i) and (v), Propositions 5.5 and 5.8] by using induction on n . □

Lemma 4.3. ([1], Proposition 5.8) *Let M be a connected matroid and $M/p = M_1 \oplus M_2$, where both M_1 and M_2 are nonempty. Then $M = (M - S_{M_1}) \parallel^p (M - S_{M_2})$.*

Lemma 4.4. ([1], Theorem 6.16(i), (v) and Propositions 4.6, 4.7) *Let $M = M_1 *^p M_2$.*

- (a) $r(M) = r(M_1) + r(M_2)$.
- (b) $c(M) = c(M_1) + c(M_2) - 1$.
- (c) $M - p = (M_1 - p) \oplus (M_2 - p)$.
- (d)

$$P(M, t) = (t-2)(t-1)^{-1}P(M_1, t)P(M_2, t) + P(M_1, t)P(M_2/p, t) + P(M_1/p, t)P(M_2, t).$$

Lemma 4.5. ([1], Proposition 4.10) *Let M be a connected matroid and $M - p = M_1 \oplus M_2$, where both M_1 and M_2 are nonempty. Then*

$$M = (M/S_{M_1}) *^p (M/S_{M_2}).$$

Lemma 4.6. *Let $M = M_1 \oplus_2^p M_2$. Then*

- (a) $r(M) = r(M_1) + r(M_2) - 1$
- (b) $c(M) = c(M_1) + c(M_2) - 1$
- (c)

$$P(M, t) = P(M_1/p, t)P(M_2 - p, t) + P(M_2, t)[(t-1)^{-1}P(M_1, t) - P(M_1/p, t)].$$

Proof. We obtain (a) and (b) by applying Lemma 2.1(c) to Lemma 4.2(a) and (b). To prove (c) we let $M' = M_1 \parallel^p M_2$. By Lemma 3.2,

$$P(M, t) = P(M', t) + P(M'/p, t). \quad (4.1)$$

By Lemma 4.2(d),

$$P(M', t) = P(M_1, t)P(M_2, t)(t-1)^{-1}. \quad (4.2)$$

By Lemma 4.2(c), $M'/p = M_1/p \oplus M_2/p$, so by Lemma 3.4,

$$P(M'/p, t) = P(M_1/p, t)P(M_2/p, t). \quad (4.3)$$

Hence, by equations (4.1), (4.2) and (4.3),

$$P(M, t) = P(M_1, t)P(M_2, t)(t-1)^{-1} + P(M_1/p, t)P(M_2/p, t). \quad (4.4)$$

By Lemma 3.2, $P(M_2/p, t) = P(M_2 - p, t) - P(M_2, t)$, so from (4.4) we obtain

$$\begin{aligned} (t-1)P(M, t) &= P(M_1, t)P(M_2, t) \\ &\quad + (t-1)P(M_1/p, t)(P(M_2 - p, t) - P(M_2, t)) \\ &= (t-1)P(M_1/p, t)P(M_2 - p, t) \\ &\quad + P(M_2, t)(P(M_1, t) - (t-1)P(M_1/p, t)). \quad \square \end{aligned}$$

Lemma 4.7. *Let M be a matroid that can be expressed as a 2-sum about a basepoint p . Then there exist matroids M_1, M_2, \dots, M_k such that $M = M_1 \oplus_2^p M_2 \oplus_2^p \cdots \oplus_2^p M_k$ and M_i/p is connected for $1 \leq i \leq k$.*

Proof. Suppose the number of components of the 2-sum, k , is as large as possible. Suppose the lemma is false and there exists a component M_i , $1 \leq i \leq k$ such that M_i/p is disconnected. Without loss of generality assume this component is M_k . Because M_k/p is disconnected, there exist matroids N_k and N_{k+1} such that $M_k/p = N_k \oplus N_{k+1}$. By Lemma 4.3, M_k is the parallel connection of some matroids N'_k and N'_{k+1} about basepoint p . But this means we have $M = M_1 \oplus_2^p M_2 \oplus_2^p \cdots \oplus_2^p M_{k-1} \oplus_2^p N_k \oplus_2^p N_{k+1}$, contradicting the maximality of k and completing the proof of Lemma 4.7. \square

Lemma 4.8. *Let $M = M_1 \oplus_2^p M_2$ and $q \in S_{M_1} - \{p\}$, where q is not parallel to p in M_1 . Then $(M_1 \oplus_2^p M_2)/q = (M_1/q \oplus_2^p M_2)$.*

Proof. This follows from [1, Proposition 5.8], the definition of 2-sum, and Lemma 2.4. \square

Lemma 4.9. *Let M_1, M_2 and M_3 be matroids such that $S_{M_i} \cap S_{M_j} = \{p\}$ for $1 \leq i < j \leq 3$. Then $M_1 \oplus_2^p (M_2 *^p M_3) = (M_1 *^p M_2) \oplus_2^p M_3$.*

Proof. It can be seen that the set of circuits of the matroids on each side of the inequality is given by $\mathcal{C}_{M_1-p} \cup \mathcal{C}_{M_2-p} \cup \mathcal{C}_{M_3-p} \cup \{(C_1-p) \cup (C_2-p) \cup (C_3-p) : p \in C_i \in \mathcal{C}_{M_i}, 1 \leq i \leq 3\}$. \square

Lemma 4.10. ([7], 2.6) *A connected matroid M is not 3-connected if and only if $M = M_1 \oplus_2^p M_2$ for some matroids M_1 and M_2 .*

5. Generalized coloops and 3-circuits.

Let $M_\Delta = A \parallel^p B$, where A and B are each isomorphic to $U_{2,3}$, the graphic matroid of a circuit of length three. Thus M_Δ is a matroid with five elements, one of which is distinguished as p , such that p is in a pair of 3-circuits, and such that $M_\Delta - p$ is a 4-circuit. Because A and B both have rank 2, it follows from Lemma 4.2(a) that the rank of M_Δ is 3. As A and B are both connected and graphic, by Lemmas 4.2(b) and 4.1, it follows that M_Δ is also connected and graphic.

Given a matroid M we define a *parallel subdivision* of M to be a matroid M_{PS} obtained by first choosing a copy M_Δ^i of M_Δ with distinguished element p_i , then relabelling one of the elements of M as p_i , and putting $M_{PS} = M \oplus_2^{p_i} M_\Delta^i$. By Lemma 4.6(a), $r(M_{PS}) = r(M) + r(M_\Delta) - 1 = r(M) + 2$. By Lemma 4.6(b), if M is connected so is M_{PS} . It also follows, by Lemma 4.1, that if M is graphic so is M_{PS} . This operation has the effect of replacing $p_i \in M$ by a pair of 2-cocircuits, which themselves form a 4-circuit.

A *generalized p -coloop* is either the rank one matroid with one element p , or $M_\Delta - p$, or any matroid that can be obtained from $M_\Delta - p$ by a sequence of parallel subdivisions. Given a generalized p -coloop (other than the one-element matroid) $M = (M_\Delta - p) \oplus_2^{p_i}$

$M_{\Delta}^1 \oplus_2^{p_2} M_{\Delta}^2 \oplus_2^{p_3} \cdots \oplus_2^{p_n} M_{\Delta}^n$, we define its p -extension to be the matroid $M' = M_{\Delta} \oplus_2^{p_1} M_{\Delta}^1 \oplus_2^{p_2} M_{\Delta}^2 \oplus_2^{p_3} \cdots \oplus_2^{p_n} M_{\Delta}^n$. Thus the p -extension of $M_{\Delta} - p$ is M_{Δ} and, in general, $p \in S_{M'}$.

Because $M_{\Delta} - p$ has rank 3, is connected and is graphic, by Lemma 4.6(b) and Lemma 4.1 its parallel subdivisions are connected and graphic. Hence generalized p -coloops are the graphic matroids corresponding to the generalized edges defined in [4]. Because every parallel subdivision increases the rank by 2, it follows that all generalized p -coloops have odd rank.

Lemma 5.1. *Let M be a generalized p -coloop with at least four elements and M' be the p -extension of M . Then M'/p is disconnected with exactly two components.*

Proof. Since M_{Δ}/p is disconnected with exactly two components we may suppose that $|M| > 4$. Let $M = (M_{\Delta} - p) \oplus_2^{p_1} M_{\Delta}^1 \oplus_2^{p_2} M_{\Delta}^2 \oplus_2^{p_3} \cdots \oplus_2^{p_n} M_{\Delta}^n$. Then $M' = M_{\Delta} \oplus_2^{p_1} M_{\Delta}^1 \oplus_2^{p_2} M_{\Delta}^2 \oplus_2^{p_3} \cdots \oplus_2^{p_n} M_{\Delta}^n$ and $M'/p = (M_{\Delta} \oplus_2^{p_1} M_{\Delta}^1 \oplus_2^{p_2} M_{\Delta}^2 \oplus_2^{p_3} \cdots \oplus_2^{p_n} M_{\Delta}^n)/p = (M_{\Delta}/p) \oplus_2^{p_1} M_{\Delta}^1 \oplus_2^{p_2} M_{\Delta}^2 \oplus_2^{p_3} \cdots \oplus_2^{p_n} M_{\Delta}^n$, by Lemma 4.8. Since M_{Δ}/p has exactly two components, it follows by repeated application of Lemma 4.6(b) that M'/p is disconnected with exactly two components. \square

A *generalized 3-circuit* is either the graphic matroid of a 3-circuit, $U_{2,3}$, or any matroid which can be obtained from $U_{2,3}$ by a sequence of parallel subdivisions.

Because $U_{2,3}$ has rank 2, is connected, and is graphic, it follows from Lemma 4.6(b) and Lemma 4.1 that all generalized 3-circuits have even rank, are connected and graphic. Hence generalized 3-circuits are the graphic matroids corresponding to the generalized triangles defined in [4].

Lemma 5.2. *Let M be a matroid with the following properties.*

- (a) M is connected.
- (b) For every $e \in S$, $M - e$ is disconnected with exactly two components.
- (c) Whenever M is the 2-sum of n connected matroids M_i ($1 \leq i \leq n$), with basepoint p , such that M_i/p is connected for all i , $1 \leq i \leq n$, then n is odd.
- (d) Whenever M is the 2-sum of two matroids M'_1 and M'_2 , with basepoint q and such that $M_1 = M'_1 - q$ is a generalized q -coloop, then $M_2 = M'_2 - q$ has exactly two components. Then M is a generalized 3-circuit.

Proof. Suppose Lemma 5.2 is false and let M be a counterexample with $|M|$ as small as possible. Because M is connected and, by (b), has no parallel elements, if $|M| = 3$ then $M = U_{2,3}$ and hence is a generalized 3-circuit. So it follows that $|M| \geq 4$.

Claim 1. M can be expressed as a 2-sum of matroids Y_1 , Y_2 and Y_3 about the same basepoint p .

Proof. By (a) and (b), M is connected but not 3-connected. Thus, by Lemma 4.10, M is the 2-sum of matroids, M'_1 and M'_2 say, about basepoint p , $p \notin M$. Claim 1 now follows from Lemma 4.7 and (c). \square

Claim 2. M can be expressed as a 2-sum of matroids Y_1 , Y_2 and Y_3 about some basepoint p , where Y_1 and Y_2 are both 3-circuits.

Proof. By Claim 1, M can be expressed as $M = Y_1 \oplus_2^p Y_2 \oplus_2^p Y_3$. Suppose p , Y_1 and Y_2 have been chosen such that $|Y_1 \oplus_2^p Y_2|$ is as small as possible. It follows from (a) and Lemma 4.6(b) that Y_1 , Y_2 and Y_3 are connected. Assume that Y_1 is not 3-connected. Hence, by Lemma 4.10, Y_1 can be expressed as a 2-sum. By Lemma 4.7, $Y_1 = X_1 \oplus_2^q X_2 \oplus_2^q \cdots \oplus_2^q X_k$, and X_i/q is connected, ($1 \leq i \leq k$). Using the minimality of $|Y_1 \oplus_2^p Y_2|$, we have $q \neq p$. Without loss of generality we may assume $p \in X_k$. We now have

$$\begin{aligned} M &= (X_1 \oplus_2^q X_2 \oplus_2^q \cdots \oplus_2^q X_{k-1} \oplus_2^q X_k) \oplus_2^p Y_2 \oplus_2^p Y_3 \\ &= X_1 \oplus_2^q X_2 \oplus_2^q \cdots \oplus_2^q X_{k-1} \oplus_2^q (X_k \oplus_2^p Y_2 \oplus_2^p Y_3). \end{aligned}$$

As X_k/q is connected and $(X_k \oplus_2^p Y_2 \oplus_2^p Y_3)/q = (X_k/q \oplus_2^p Y_2 \oplus_2^p Y_3)$ by Lemma 4.8, it follows from Lemma 4.6(b) that $(X_k \oplus_2^p Y_2 \oplus_2^p Y_3)/q$ is connected. Thus k is odd, by (c), and $k \geq 3$. Now $M = X_1 \oplus_2^q X_2 \oplus_2^q (X_3 \oplus_2^q \cdots \oplus_2^q X_{k-1} \oplus_2^q X_k \oplus_2^p Y_2 \oplus_2^p Y_3)$. Since $X_1 \oplus_2^q X_2$ is contained in Y_1 , this contradicts the minimality of $|Y_1 \oplus_2^p Y_2|$. Hence Y_1 is 3-connected. By a similar argument, Y_2 is also 3-connected.

Because M is a 2-sum, $|Y_1|, |Y_2| \geq 3$. Also, every element of $Y_1 - p$ is in M , so we can choose $f \in M \cap Y_1$. But, by Lemma 2.6, provided $|Y_1| \geq 4$, $Y_1 - f$ is connected and hence $M - f$ is connected, contradicting (b). Hence $|Y_1| = 3$, and by a similar argument, $|Y_2| = 3$. Since M has no parallel elements, it follows that Y_1 and Y_2 are 3-circuits, so completing the proof of Claim 2. \square

Claim 3. M is a 2-sum of matroids M'_1 and M'_2 , with basepoint p such that $M'_1 = M_\Delta$ and so $M_1 = M'_1 - p$ is a generalized p -coloop.

Proof. Let Y_1 , Y_2 and Y_3 be as in Claim 2. Let M'_1 be the parallel connection of Y_1 and Y_2 and $M'_2 = Y_3$. Then $M = M'_1 \oplus_2^p M'_2$ and $M'_1 = M_\Delta$. Hence $M_1 = M'_1 - p$ is a generalized p -coloop. \square

We now return to the proof of Lemma 5.2. Let p , M'_1 and M'_2 be defined as in Claim 3. Thus M_1 is a generalized p -coloop.

Consider M'_2 . We will show that it satisfies the hypotheses of Lemma 5.2.

- (a) Because M is connected, by Lemma 4.6(b), M'_2 is also connected.
- (b) Choose $e \in M'_2$. If $e \neq p$, then $e \in S_M$. Applying (b) to M , we deduce that $M - e$ is disconnected with exactly two components. Since $M - e = M_\Delta \oplus_2^p (M'_2 - e)$, it follows from Lemma 4.6(b) that $M'_2 - e$ is disconnected with exactly two components. If $e = p$, then $M'_2 - e = M_2$, which has exactly two components since M satisfies (d).
- (c) Choose matroids N'_i ($1 \leq i \leq n$), with basepoint q , such that N'_i/q is connected for all i , $1 \leq i \leq n$ and M'_2 is their 2-sum. Hence $q \notin M'_2$, so $q \neq p$. Assume without loss of generality that $p \in N'_1$. Because M'_1 and N'_1 are both connected, $L'_1 = M'_1 \oplus_2^p N'_1$ is connected and $q \in L'_1$. Similarly, $L'_i/q = M'_1 \oplus_2^p N'_i/q$ is connected. But $M = L'_1 \oplus_2^q N'_2 \oplus_2^q \cdots \oplus_2^q N'_n$, and because M satisfies (c), n is odd. Thus M'_2 also satisfies (c).

(d) Let M'_2 be the 2-sum of matroids N'_1 and N'_2 with basepoint q such that $N_1 = N'_1 - q$ is a generalized q -coloop. We need to show that $N_2 = N'_2 - q$ has exactly two components. Because M satisfies (d) we know $M_2 = M'_2 - p$ has exactly two components and $p \in M'_2$, so $p \neq q$. There are two cases to consider.

- Suppose $p \in N'_2$. Then $M = N'_1 \oplus_2^q L'_2$, where $L'_2 = N'_2 \oplus_2^p M_\Delta$. Since N_1 is a generalized q -coloop we may apply (d) to M and deduce that $L_2 = L'_2 - q$ has exactly two components. But $L_2 = N_2 \oplus_2^p M_\Delta$, so it follows from Lemma 4.6(b) that N_2 also has exactly two components.
- Suppose $p \in N'_1$. Then $M = L'_1 \oplus_2^q N'_2$, where $L'_1 = N'_1 \oplus_2^p M_\Delta$. Because N_1 is a generalized q -coloop so is L_1 . Applying (d) to M we deduce that N_2 has exactly two components.

Since M'_2 satisfies all four conditions and has fewer elements than M , it follows that M'_2 is a generalized 3-circuit. Since $M = M'_2 \oplus_2^p M_\Delta$, we deduce that M is also a generalized 3-circuit. This contradicts the choice of M and completes the proof of Lemma 5.2. \square

6. Zeros of characteristic polynomials

Theorem 6.1. *Let M be a loopless matroid with rank r and characteristic polynomial $P(M, t)$. Then*

- (a) $P(M, t) = t^r - |M|t^{r-1} + k_{r-2}t^{r-2} - \cdots + (-1)^r k_0$ where k_0, \dots, k_{r-2} are positive integers
- (b) $(-1)^r P(M, t) \geq (1-t)^r$ for $t \in (-\infty, 1)$
- (c) $P(M, 1) = 0$.

Proof. We use induction on $|M|$. If M is a single element matroid then $P(M, t) = t - 1$, which satisfies (a),(b) and (c). Using Lemma 3.5 we may assume that M has no parallel elements. When $|M| \geq 2$ we choose $e \in S$ and use Lemmas 2.1, 2.3, 3.2 and the inductive hypothesis on $M - e$ and M/e if e is not a coloop, and use Lemmas 2.1, 3.3 and the inductive hypothesis on $M - e$ if e is a coloop. \square

Theorem 6.2. *Let M be a loopless matroid. Then the multiplicity of 1 as a zero of $P(M, t)$ is equal to the number of components of M .*

Proof. We adopt the proof technique of Woodall [12] for chromatic polynomials of graphs. If M is a single element matroid, $P(M, t) = t - 1$, and the theorem holds. Hence, suppose $|M| \geq 2$. Using Lemma 3.5 we may assume that M has no parallel elements.

Suppose M is connected. We show that $\frac{d}{dt}P(M, t)$ is nonzero at $t = 1$. Choose $e \in S$. By Lemma 3.2,

$$\frac{d}{dt}P(M, t) = \frac{d}{dt}P(M - e, t) - \frac{d}{dt}P(M/e, t). \quad (6.1)$$

Using Theorem 6.1(b) and (c) and Lemma 2.1(c), $\frac{d}{dt}P(M - e, t)$ has sign $(-1)^{r(M)-1}$ or is zero when $t = 1$ and $\frac{d}{dt}P(M/e, t)$ has sign $(-1)^{r(M)}$ or is zero when $t = 1$. Thus $\frac{d}{dt}P(M - e, t)$ and $\frac{d}{dt}P(M/e, t)$ are either zero or have opposite signs at $t = 1$. By Lemma 2.5, either

$M - e$ or M/e is connected. Thus, by induction, either $\frac{d}{dt}P(M - e, t)$ or $\frac{d}{dt}P(M/e, t)$ is nonzero at $t = 1$. Hence, by (6.1), we have $\frac{d}{dt}P(M, t)$ is nonzero at $t = 1$.

Finally, if M is not connected then we apply the inductive hypothesis to each connected component of M and use Lemma 3.4. \square

Theorem 6.3. *Let M be a loopless matroid with rank $r(M)$, and $c(M)$ components. Then*

$$(-1)^{r(M)+c(M)}P(M, t) \geq (t - 1)^{r(M)} \tag{6.2}$$

for $t \in (1, \frac{32}{27}]$.

Proof. Suppose Theorem 6.3 is false. Let M be a matroid and $t \in (1, \frac{32}{27}]$ such that $P(M, t)$ does not satisfy (6.2), and assume that $|M|$ is as small as possible. Clearly $|M| \geq 2$. Using Lemma 3.5 we may deduce that M has no parallel elements. We shall show that M satisfies the hypotheses of Lemma 5.2 and hence is a generalized 3-circuit.

Claim 4. *M is connected.*

Proof. We proceed by contradiction. Suppose $M = M_1 \oplus M_2$. By Lemma 3.4, $P(M, t) = P(M_1, t)P(M_2, t)$. As M is the smallest counterexample, M_1 and M_2 satisfy (6.2). Thus

$$\begin{aligned} (-1)^{r(M)+c(M)}P(M, t) &= (-1)^{r(M_1)+c(M_1)}P(M_1, t)(-1)^{r(M_2)+c(M_2)}P(M_2, t) \\ &\geq (t - 1)^{r(M_1)+r(M_2)} \\ &= (t - 1)^{r(M)}. \end{aligned}$$

This gives the required contradiction. \square

Claim 5. *M/e is connected for all $e \in S$.*

Proof. Suppose M/e is not connected. By Lemmas 4.3 and Claim 4, M is the parallel connection of two connected matroids M_1 and M_2 , about basepoint e . So, by Lemma 4.2, we have $r(M) = r(M_1) + r(M_2) - 1$ and

$$P(M, t) = P(M_1, t)P(M_2, t)/(t - 1).$$

Applying (6.2) to M_1 and M_2 we obtain

$$\begin{aligned} (-1)^{r(M)+1}P(M, t) &= (-1)^{r(M_1)+1}P(M_1, t)(-1)^{r(M_2)+1}P(M_2, t)/(t - 1) \\ &\geq (t - 1)^{r(M_1)+r(M_2)}/(t - 1) \\ &= (t - 1)^{r(M)}. \end{aligned}$$

This gives the required contradiction. \square

Claim 6. *$M - e$ has exactly two components for all $e \in S$.*

Proof. The proof will be in two stages. We first show that $M - e$ has an even number of components for all $e \in S$, and hence $M - e$ is disconnected.

Using Lemmas 2.1(c), 2.3(b), 3.2 and Claim 4 we deduce that

$$(-1)^{r(M)+1}P(M, t) = (-1)^{r(M-e)+1}P(M-e, t) + (-1)^{r(M/e)+1}P(M/e, t).$$

Thus, if $c(M-e)$ is odd we have

$$(-1)^{r(M)+1}P(M, t) = (-1)^{r(M-e)+c(M-e)}P(M-e, t) + (-1)^{r(M/e)+1}P(M/e, t).$$

Using Claim 5 and applying (6.2) to $M-e$ and M/e gives

$$\begin{aligned} (-1)^{r(M)+1}P(M, t) &\geq (t-1)^{r(M-e)} + (t-1)^{r(M/e)} \\ &= (t-1)^{r(M)} + (t-1)^{r(M)-1}. \end{aligned}$$

This contradicts the choice of M and t and hence $c(M-e)$ is even.

Let $M-e = M_1 \oplus \cdots \oplus M_{2r}$, $X_1 = M_1 \oplus \cdots \oplus M_r$, $X_2 = M_{r+1} \oplus \cdots \oplus M_{2r}$. Then $M-e = X_1 \oplus X_2$. By Lemmas 4.5 and 4.4(a),(b), $M = N_1 *^c N_2$, where $N_1 = M/S_{X_1}$, and $N_1 = M/S_{X_1}$, $N_2 = M/S_{X_2}$ are connected, and $r(M) = r(N_1) + r(N_2)$. By Lemma 4.4(d),

$$\begin{aligned} (-1)^{r(M)+1}P(M, t) &= -\left(\frac{t-2}{t-1}\right) (-1)^{r(N_1)+1}P(N_1, t) (-1)^{r(N_2)+1}P(N_2, t) \\ &\quad + (-1)^{r(N_1)+1}P(N_1, t) (-1)^{r(N_2/e)+1}P(N_2/e, t) \\ &\quad + (-1)^{r(N_1/e)+1}P(N_1/e, t) (-1)^{r(N_2)+1}P(N_2, t). \end{aligned}$$

If N_1/e and N_2/e are both connected then, applying (6.2), we obtain

$$\begin{aligned} (-1)^{r(M)+1}P(M, t) &\geq \left(\frac{2-t}{t-1}\right) (t-1)^{r(N_1)+r(N_2)} \\ &\quad + (t-1)^{r(N_1)+r(N_2/e)} + (t-1)^{r(N_1/e)+r(N_2)} \\ &= (4-t)(t-1)^{r(M)-1} \\ &> (t-1)^{r(M)} \end{aligned}$$

since $t \in (1, \frac{32}{27}]$. This contradiction implies that at least one of N_1/e and N_2/e is disconnected. Without loss of generality, N_1/e is disconnected. Then, by Lemma 2.5, N_1-e is connected. However, using Lemma 2.4, we have

$$N_1-e = M/S_{X_1} - e = (M-e)/S_{X_1} = (X_1 \oplus X_2)/S_{X_1} = X_2.$$

Thus, $X_2 = M_{r+1} \oplus \cdots \oplus M_{2r}$ is connected. Hence $r = 1$ and $M-e$ has exactly two components. \square

Claim 7. Suppose M is the 2-sum of n connected matroids M_i , with basepoint p , such that M_i/p is connected for all i , $1 \leq i \leq n$. Then n must be odd.

Proof. Let M' denote the parallel connection of M_1, M_2, \dots, M_n . Hence $M = M' - p$, and by Lemma 4.2(c), $M'/p = M_1/p \oplus M_2/p \oplus \cdots \oplus M_n/p$. By Lemmas 3.2, 4.2(d), 3.4,

$$P(M, t) = P(M', t) + P(M'/p, t),$$

$$P(M', t) = (t-1)^{-n+1} \prod_{i=1}^n P(M_i, t)$$

and

$$P(M'/p, t) = \prod_{i=1}^n P(M_i/p, t).$$

Hence

$$P(M, t) = (t - 1)^{-n+1} \prod_{i=1}^n P(M_i, t) + \prod_{i=1}^n P(M_i/p, t).$$

Using Lemmas 4.2(a) and 2.3(b) gives

$$\begin{aligned} (-1)^{r(M)+1} P(M, t) &= (t - 1)^{-n+1} \prod_{i=1}^n (-1)^{r(M_i)+1} P(M_i, t) \\ &\quad + (-1)^n \prod_{i=1}^n (-1)^{r(M_i/p)+1} P(M_i/p, t). \end{aligned}$$

If n is even then we may apply (6.2) to $P(M_i, t)$ and $P(M_i/p, t)$ for $1 \leq i \leq n$, to give

$$\begin{aligned} (-1)^{r(M)+1} P(M, t) &\geq (t - 1)^{-n+1} \prod_{i=1}^n (t - 1)^{r(M_i)} + \prod_{i=1}^n (t - 1)^{r(M_i/p)} \\ &= (t - 1)^{r(M)} + (t - 1)^{r(M)-1}. \end{aligned}$$

This contradicts the choice of M and hence n must be odd. □

Claim 8. Suppose $M = M_1 \oplus_2^q M_2$ where $M_1 - q$ is a generalized q -coloop. Then $M_2 - q$ has exactly two components.

Proof. Suppose $|M_2| \leq 3$. Then $|M_2| = 3$ by the definition of 2-sum and, since M_2 is connected and has no parallel elements, $M_2 = U_{2,3}$. Thus the claim holds. Henceforth we will assume that $|M_2| > 3$.

Claim 8(a). $(t - 1)^{-1} P(M_1, t) - P(M_1/q, t) > 0$.

Proof. Let $N = M_1 \oplus_2^q U_{2,3}$. By Lemma 4.6(b), N is connected. Since $M_1 - q$ is a generalized q -coloop we have that $r(M_1 - q)$ is odd. By Lemma 2.1(c), $r(M_1) = r(M_1 - q)$. Using Lemma 4.6(a) and the fact that $r(U_{2,3}) = 2$, we have $r(N) = r(M_1) + 1$. Thus $r(N)$ is even.

Applying (6.2) inductively to N we deduce that

$$P(N, t) < 0. \tag{6.3}$$

Applying Lemma 4.6(c) to N with $N = M_1 \oplus_2^q U_{2,3}$, we obtain

$$\begin{aligned} P(N, t) &= P(M_1/q, t)(t - 1)^2 \\ &\quad + (t - 1)(t - 2)[(t - 1)^{-1} P(M_1, t) - P(M_1/q, t)]. \end{aligned} \tag{6.4}$$

Since $r(M_1)$ is odd, it follows from Lemma 2.3(b) that $r(M_1/q)$ is even. Since $M_1 - q$ is a generalized q -coloop, by Lemma 5.1, M_1/q is disconnected with exactly two components. Applying (6.2) to M_1/q we deduce that

$$P(M_1/q, t) > 0. \tag{6.5}$$

Since $t \in (1, \frac{32}{27}]$, it follows from (6.3) that $(t-1)P(N, t) < 0$, from (6.5) that $(t-1)^2P(M_1/q, t) > 0$, and that $(t-1)(t-2) < 0$. Hence, from (6.4), we obtain

$$(t-1)^{-1}P(M_1, t) - P(M_1/q, t) > 0. \quad \square$$

Claim 8(b). $c(M_2 - q)$ is even.

Proof. By Lemma 4.6 and Claim 4, M_1 and M_2 are connected, and

$$(-1)^{r(M)+1}P(M, t) = (-1)^{r(M_1/q)+1}P(M_1/q, t)(-1)^{r(M_2-q)}P(M_2 - q, t) + \alpha\beta,$$

where $\alpha = (-1)^{r(M_2)+1}P(M_2, t)$ and

$$\beta = (t-1)^{-1}(-1)^{r(M_1)+1}P(M_1, t) - (-1)^{r(M_1/q)}P(M_1/q, t).$$

Applying (6.2) to M_2 gives $\alpha > 0$. Using Claim 8(a) and the fact that $r(M_1)$ is odd we may deduce that $\beta > 0$. Thus

$$(-1)^{r(M)+1}P(M, t) \geq (-1)^{r(M_1/q)}P(M_1/q, t)(-1)^{r(M_2-q)+1}P(M_2 - q, t).$$

If $c(M_2 - q)$ is odd then, since $c(M_1)/q$ has exactly two components, we may apply (6.2) to deduce that

$$\begin{aligned} (-1)^{r(M)+1}P(M, t) &\geq (t-1)^{r(M_1/q)+r(M_2-q)} \\ &= (t-1)^{r(M)}. \end{aligned}$$

This contradiction implies that $c(M_2 - q)$ is even. \square

Let $M_2 - q = X_1 \oplus \cdots \oplus X_{2r}$, $Y_1 = X_1 \oplus \cdots \oplus X_r$ and $Y_2 = X_{r+1} \oplus \cdots \oplus X_{2r}$. Then $M_2 - q = Y_1 \oplus Y_2$ so, by Lemma 4.5, $M_2 = N_1 *^q N_2$ where $N_1 = M_2/S_{Y_2}$ and $N_2 = M_2/S_{Y_1}$. Thus $M = M_1 \oplus_2^q (N_1 *^q N_2)$. Since M is connected, it follows from Lemmas 4.6(b) and 4.4(b) that M_1 , N_1 and N_2 are connected. By Lemma 4.9, $M = (M_1 *^q N_1) \oplus_2^q N_2$.

Suppose $r \geq 2$. Then, since $N_1 - q = Y_1$, $N_1 - q$ is disconnected so, by Lemma 2.5, N_1/q is connected. Since M_1 and N_1 are connected, it follows from Lemma 4.6(b) that $(M_1 *^q N_1)/q = M_1 \oplus_2^q N_1$ is connected. Now the fact that $M = (M_1 *^q N_1) \oplus_2^q N_2$ contradicts Claim 7. Thus $r = 1$. \square

Using Claims 4, 6, 7 and 8 it follows that M satisfies the hypotheses of Lemma 5.2 and hence M is a generalized 3-circuit. Thus M is graphic. But, by [4, Theorem 5] and [8, Theorem 2.4], Theorem 6.3 is known to be true for graphic matroids. This contradicts the choice of M as a counterexample, and completes the proof of Theorem 6.3. \square

Lemma 6.1. ([8], Theorem 2.5) *Let $t_0 > \frac{32}{27}$, $\epsilon > 0$. Then there exists a graph G such that $P(G, t)$ has a zero in $(t_0 - \epsilon, t_0 + \epsilon)$.*

Corollary 6.1. *The maximal zero-free intervals for characteristic polynomials of loopless matroids are precisely $(-\infty, 1)$ and $(1, \frac{32}{27}]$.*

Proof. The result follows directly from Theorem 6.1, Theorem 6.3 and Lemma 6.1. \square

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